Non-commutative Wick polynomials



1. Motivation

- 2. Moments and cumulants
- 3. Free Wick polynomials
- 4. If time permits:
 - 4.1 Modification of products
 - 4.2 Relation to power series



Definition

Let X be a r.v. with $\mathbb{E}X^n < \infty$ for all n > 0. Recursive definition:

$$W'_n(x) = nW_{n-1}(x), \quad \mathbb{E}W_n(X) = 0.$$

For example:
$$W_1(x) = x - \mathbb{E}X$$
, $W_2(x) = x^2 - 2x\mathbb{E}X + 2(\mathbb{E}X)^2 - \mathbb{E}X^2$,...

Definition (Multivariate Wick polynomials)

$$\frac{\partial}{\partial x_i}W_n(x_1,\ldots,x_n)=W_{n-1}(x_1,\ldots,x_{i-1},x_{i+1},\ldots,x_n),\quad \mathbb{E}W_n(X_1,\ldots,X_n)=0.$$



Let $\mathcal{F}^{\circ} \coloneqq \mathbb{C}\Omega \oplus H \oplus H^{\circ 2} \oplus \cdots$ be the symmetric Fock space over *H*.

For each $f \in H$ we have (bosonic) annihilation and creation operators a(f), $a^{\dagger}(f)$ on \mathcal{F}° such that

$$a(f)\Omega = 0, \quad a^{\dagger}(f)\Omega = f$$

and

$$a(f)(f_1 \circ \cdots \circ f_n) = \sum_{j=1}^n \langle f, f_j \rangle f_1 \circ \cdots \circ f_{j-1} \circ f_{j+1} \circ \cdots \circ f_n,$$

$$a^{\dagger}(f)(f_1 \circ \cdots \circ f_n) = f \circ f_1 \circ \cdots \circ f_n.$$

They satisfy the (canonical) commutation relation

$$a(f)a^{\dagger}(g) - a^{\dagger}(g)a(f) = \langle f, g \rangle 1.$$



The normal order operator N puts creation operators to the left of annihilation operators.

For example $N(a^{\dagger}(f)a(g)) = a^{\dagger}(f)a(g)$ and $N(a(f)a^{\dagger}(g)) = a^{\dagger}(g)a(f)$, etc.

The (unnormalized) position operators $p(f) \coloneqq a(f) + a^{\dagger}(f)$ satisfy N(p(f)) = p(f) and

$$\begin{split} p(f)p(g) &= a(f)a(g) + a(f)a^{\dagger}(g) + a^{\dagger}(f)a(g) + a^{\dagger}(f)a^{\dagger}(g) \\ &= a(f)a(g) + a^{\dagger}(g)a(f) + a^{\dagger}(f)a(g) + a^{\dagger}(f)a^{\dagger}(g) + \langle f, g \rangle 1 \\ &= \mathsf{N}(p(f)p(g)) + \langle f, g \rangle 1, \end{split}$$

i.e. $N(p(f)p(g)) = p(f)p(g) - \langle f, g \rangle 1$.

Denoting X = p(f), Y = p(g) and $\mathbb{E}(b) \coloneqq \langle b\Omega, \Omega \rangle$ we see that $N(XY) = XY - \mathbb{E}(XY) = W_2(X, Y)$. By definition $\mathbb{E}[N(p(f_1) \cdots p(f_n))] = 0$.



Suppose we have a measure of the form $\mu(dx) = e^{-S(x)} dx$ and set $I[f] = \int f d\mu$.

Integrating by parts we get, for any nice function f, the Dyson-Schwinger equation

$$I[\partial_i f - (\partial_i S)f] = 0.$$

The measure μ is characterized by the values I[f], for nice f, usually polynomials are enough.

For a 1-D Gaussian weight the equation is simply I(f' - xf) = 0 which entails $I(x^{2n}) = (n - 1)!!$ and zero else.

This ultimately means that $Z[J] \coloneqq I_x(e^{Jx}) = e^{\frac{1}{2}J^2}$, or $I[e^{Jx-\frac{1}{2}J^2}] = 1$. Observe also that $(J - \frac{d}{dJ})Z[J] = 0$.

Actually, the Dyson-Schwinger equaition implies that $I(H_{n+1}) = -I(H'_n - xH_n) = 0$, so that re-expanding x^n in the H_n basis:

$$I(x^n) = \sum_{k=0}^n \alpha_k I(H_k) = \alpha_0$$



More generally, for a multidimensional Gaussian weight $S(x) = \frac{1}{2}g_{ij}x^ix^j$ a similar argument gives that

$$I(\mathbf{x}^{j}\mathbf{x}^{i_{1}}\cdots\mathbf{x}^{i_{n}})=\sum_{k=1}^{n}\mathbf{g}^{ij_{k}}I(\mathbf{x}^{i_{1}}\cdots\hat{\mathbf{x}}^{i_{k}}\cdots\mathbf{x}^{i_{n}}).$$

But, since we have that $I(H_n) = 0$ we can again re-expand any monomial in the H_n basis and recover the above formula from our knowledge of H_n .

In our annihilation-creator operator example:

$$\langle p(f_1)p(f_2)\Omega, \Omega \rangle = \langle f_1, f_2 \rangle \langle p(f_1)p(f_2)p(f_3)p(f_4)\Omega, \Omega \rangle = \langle f_1, f_2 \rangle \langle f_3, f_4 \rangle + \langle f_1, f_3 \rangle \langle f_2, f_4 \rangle + \langle f_1, f_4 \rangle \langle f_2, f_3 \rangle$$



Definition

A noncommutative probability space is a tuple (A, φ) where A is an associative algebra and $\varphi : A \to k$ is unital, *i.e.* $\varphi(1_A) = 1$.

On
$$T(A) \coloneqq \bigoplus_{n>0} A^{\otimes n}$$
 define $\Delta \colon T(A) \to T(A) \otimes T(A)$ by
 $\Delta^{\sqcup \sqcup}(a_1 \cdots a_n) \coloneqq \sum_{S \subseteq [n]} a_S \otimes a_{[n] \setminus S}.$

This induces a product on $T(A)^*$:

 $\mu \sqcup \nu \coloneqq (\mu \otimes \nu) \Delta^{\sqcup \sqcup}.$



Define $\phi : T(A) \to k$ by $\phi(a_1 \cdots a_n) \coloneqq \phi(a_1 \cdot a_1 \cdots a_n)$ and extend to $\overline{T}(A) \coloneqq k 1 \oplus T(A)$ by $\phi(1) = 1$.

There is $c : \overline{T}(A) \to k$ with c(1) = 0 such that $\phi = \exp^{\Box}(c)$. In particular

$$\phi(a_1\cdots a_n)=\sum_{\pi\in P(n)}\prod_{B\in\pi}c(a_B).$$

Definition

Since ϕ is invertible, we set $W := (id \otimes \phi^{-1})\Delta^{\sqcup \sqcup}$.

Theorem

The map $W : \overline{T}(A) \to \overline{T}(A)$ is the unique linear map such that $\partial_a \circ W = W \circ \partial_a$ and $\phi \circ W = \varepsilon$. Its inverse is given by $W^{-1} = (id \otimes \phi)\Delta^{\sqcup}$.

Observe also that trivially $\varepsilon \circ W^{-1} = \phi = \exp^{\Box \Box}(c)$.



We have other notions of independence: freeness, boolean idependence, monotone independence, etc...

Each is characterised by a set of cumulants: κ , β , ρ resp.

On the double tensor algebra $\overline{T}(T(A))$ consider

$$\Delta(a_1\cdots a_n)\coloneqq \sum_{S\subseteq [n]}a_S\otimes a_{J_1^S}|\cdots|a_{J_k^S}.$$

This splits as

$$\Delta_{\prec}(a_{1}\cdots a_{n}) \coloneqq \sum_{1\in S\subseteq [n]} a_{S}\otimes a_{J_{1}^{S}}|\cdots|a_{J_{k}^{S}},$$
$$\Delta_{\succ}(a_{1}\cdots a_{n}) \coloneqq \sum_{1\notin S\subseteq [n]} a_{S}\otimes a_{J_{1}^{S}}|\cdots|a_{J_{k}^{S}}.$$



Therefore, the convolution product $\mu * v := (\mu \otimes v)\Delta$ also splits:

$$\mu < \mathbf{v} \coloneqq (\mu \otimes \mathbf{v}) \Delta_{<}, \quad \mu > \mathbf{v} \coloneqq (\mu \otimes \mathbf{v}) \Delta_{>}.$$

Consider Φ : $\overline{T}(T(A)) \rightarrow k$ the unique character extension of ϕ .

Theorem (Ebrahimi-Fard, Patras; 2014, 2017)

The cumulants κ , β , ρ are the unique infinitesimal characters of $\overline{T}(T(A))$ such that

 $\Phi = \varepsilon + \kappa < \Phi$ $= \varepsilon + \Phi > \beta$

and $\Phi = \exp_*(\rho)$.



Theorem (Speicher; 1997)

$$\varphi(a_1 \cdot_A \cdots \cdot_A a_n) = \sum_{\pi \in NC(n)} \prod_{B \in \pi} \kappa(a_B).$$

Theorem (Speicher, Woroudi; 1997)

$$\varphi(a_1 \cdot_A \cdots \cdot_A a_n) = \sum_{\pi \in Int(n)} \prod_{B \in \pi} \beta(a_B).$$

Theorem (Hasebe, Saigo; 2011)

$$\varphi(a_1 \cdot_A \cdots \cdot_A a_n) = \sum_{(\pi,\lambda) \in \mathcal{M}(n)} \frac{1}{|\pi|!} \prod_{B \in \pi} \rho(a_B)$$



We write

$$\Phi = \mathcal{E}_{\prec}(\kappa) = \mathcal{E}_{\succ}(\beta) = \exp_*(\rho).$$

Every character has an inverse for *. For Φ we have

$$\Phi^{-1} = \mathcal{E}_{>}(-\kappa) = \mathcal{E}_{<}(-\beta) = \exp_{*}(-\rho).$$

In fact, characters on $\overline{T}(T(A))$ form a group denoted by G.

Observe that $\Delta : T(A) \to T(A) \otimes \overline{T}(T(A))$, i.e. we have a coaction.

Thus, the character group *G* acts on End(T(A)).



Definition

By analogy, define $W : T(A) \rightarrow T(A)$ by

$$W \coloneqq (\mathsf{id} \otimes \Phi^{-1}) \Delta$$

Examples:

$$\begin{split} W(a) &= a - \phi(a) 1\\ W(ab) &= ab - a\phi(b) - b\phi(a) + (2\phi(a)\phi(b) - \phi(a \cdot b)) 1\\ W(abc) &= abc - \phi(c)ab - \phi(b)ac - \phi(a)bc\\ &- [\phi(b \cdot c) - 2\phi(b)\phi(c)]a + \phi(a)\phi(c)b - [\phi(a \cdot b) - 2\phi(a)\phi(b)]c\\ &- [\phi(a \cdot b \cdot c) - 2\phi(a)\phi(b \cdot c) - 2\phi(c)\phi(a \cdot b) - \phi(b)\phi(a \cdot c)\\ &+ 5\phi(a)\phi(b)\phi(c)]1 \end{split}$$



By definition

$$\Phi \circ W = (\Phi \otimes \Phi^{-1}) \Delta = \varepsilon$$

that is, $\Phi(W(a_1 \dots a_n)) = 0$ for any $a_1, \dots, a_n \in A$.

It's easy to check that W is invertible with $W^{-1} = (id \otimes \Phi)\Delta$ and so $\Phi = \varepsilon \circ W^{-1}$.

In particular

$$a_1 \cdots a_n = \sum_{S \subseteq [n]} W(a_s) \Phi(a_{J_1^S}) \cdots \Phi(a_{J_k^S}).$$

Theorem (Anshelevich, 2004)

$$W(a_1 \cdots a_n) = \sum_{S \subseteq [n]} a_S \sum_{\substack{\pi \in \operatorname{Int}([n] \setminus S) \\ \pi \cup S \in NC(n)}} (-1)^{|\pi|} \prod_{B \in \pi} \kappa(a_B).$$



Theorem

The Wick polynomials satisfy the recursion

$$W(a_1\cdots a_n)=a_1W(a_2\cdots a_n)-\sum_{j=0}^{n-1}W(a_{j+1}\cdots a_n)\kappa(a_1\cdots a_j).$$

Proof.

$$W = (\mathrm{id} \otimes \Phi^{-1})\Delta$$

= $\mathrm{id} < \Phi^{-1} + \mathrm{id} > \Phi^{-1}$
= $\mathrm{id} < \Phi^{-1} - \mathrm{id} > (\Phi^{-1} > \kappa)$
= $\mathrm{id} < \Phi^{-1} - W > \kappa.$



Now consider the *full Fock space* $\mathcal{F} = \mathbb{C}\Omega \oplus H \oplus H^{\otimes 2} \oplus \cdots$.

We have annihilation and creator operators

 $a(f)(f_1 \otimes \cdots \otimes f_n) = \langle f, f_1 \rangle f_2 \otimes \cdots \otimes f_n, \quad a^*(f)(f_1 \otimes \cdots \otimes f_n) = f \otimes f_1 \otimes \cdots \otimes f_n.$ This time they satisfy $a(f)a^*(g) = \langle f, g \rangle 1.$

Set as before $p(f) = (a(f) + a^*(f))$. We have

$$\langle p(f_1)p(f_2)\Omega, \Omega \rangle = \langle f_1, f_2 \rangle \\ \langle p(f_1)p(f_2)p(f_3)p(f_4)\Omega, \Omega \rangle = \langle f_1, f_2 \rangle \langle f_3, f_4 \rangle + \langle f_1, f_4 \rangle \langle f_2, f_3 \rangle.$$

We get a free version of Wick's theorem (Effros, Poppa; 2003).



Theorem

The Wick polynomials can be expressed in terms of boolean cumulants

$$W = (\mathsf{id} - \mathsf{id} > \beta) < \Phi^{-}$$

Proof.

Previous theorem plus the fact that $\kappa = \Phi > \beta < \Phi^{-1}$.

Definition

The Boolean Wick map is defined by

$$W' := \mathrm{id} - \mathrm{id} > \beta.$$

Therefore

$$W'(a_1\cdots a_n)=a_1\cdots a_n-\sum_{j=1}^n a_{j+1}\cdots a_n\beta(a_1\cdots a_j).$$



Theorem

Boolean Wick polynomials are centered

Proof.

 $\Phi \circ W' = \Phi - \Phi > \beta = \varepsilon$

Theorem

We have

$$a_1\cdots a_n=W'(a_1\cdots a_n)+\sum_{j=1}^{n-1}\Phi(a_1\cdots a_j)W'(a_{j+1}\cdots a_n).$$

From a previous computation $W' = W < \Phi$, that is

$$W'(a_1\cdots a_n)=\sum_{1\in S\subseteq [n]}W(a_S)\Phi(a_{J_1^S})\cdots\Phi(a_{J_k^S}).$$



Assume we have a second state $\psi : A \rightarrow k$.

Definition

Two-state cumulants are defined implicitly by

$$\varphi(a_1 \cdot_A \cdots \cdot_A a_n) = \sum_{\pi \in NC(n)} \prod_{B \in \text{Outer}(\pi)} R^{\varphi, \psi}(a_B) \prod_{B \in \text{Inner}(\pi)} \kappa^{\psi}(a_B).$$

Theorem (Ebrahimi-Fard, Patras; 2018)

 $R^{\varphi,\psi}$ is the unique infinitesimal character of $\overline{T}(T(A))$ such that

$$\Phi = \varepsilon + \Phi > (\Psi^{-1} > R^{\varphi, \psi} < \Psi).$$



Directly,

$$R^{\varphi,\psi}=\Psi > \beta^{\varphi} < \Psi^{-1}.$$

In particular,

$$R^{\varphi,\varphi} = \Phi > \beta^{\varphi} < \Phi^{-1} = \kappa^{\varphi}, \ R^{\varphi,\varepsilon} = \beta^{\varphi}.$$

Definition

The conditionally-free Wick polynomials are defined as

$$W^c \coloneqq W < (\Phi * \Psi^{-1}).$$

This means

$$W^{c} = \left(\mathsf{id} - \mathsf{id} > \Theta_{\Psi}(R^{\varphi, \psi}) \right) < \Psi^{-1}$$

where $\Theta_{\Psi}(\mu) \coloneqq \Psi^{-1} > \mu < \Psi$.



Back to the Fock space interpretation, we actually have a whole family parametrized by $q \in (-1, 1)$. The bosonic (i.e. symmetric) corresponds to q = 1, the free case corresponds to q = 0.

The next interesting case is the fermionic setting q = -1.

For any $q \in (-1, 1)$, the associated annihilation and creation operators are

$$a(f)(f_1 \otimes \cdots \otimes f_n) = \sum_{j=1}^n q^{j-1} \langle f, f_j \rangle f_1 \otimes \cdots \otimes \hat{f_j} \otimes \cdots \otimes f_n$$
$$a^*(f)(f_1 \otimes \cdots \otimes f_n) = f \otimes f_1 \otimes \cdots \otimes f_n.$$

satisfying the *q*-commutator relation

$$[a(f), a^*(g)]_q \coloneqq a(f)a^*(g) - qa^*(g)a(f) = \langle f, g \rangle 1.$$



Thank you!



Since *W* is invertible, one can induce a product on T(A) by

 $x \bullet y = W(W^{-1}(x)W^{-1}(y))$

Proposition

The • product admits the closed-form expression: for $x = a_1 \cdots a_n$, $y = a_{n+1} \cdots a_{n+m}$

$$x \bullet y = \sum_{S \subseteq [n+m]} a_S \Phi(a_{J_1^S}) \cdots \Phi(a_{J_k^S}).$$



The relations between moments and cumulants can also be encoded by power series.

In the classical case, one uses exponential generating functions:

$$\sum_{n\geq 0} m_n \frac{\lambda^n}{n!} = \exp\left(\sum_{k>0} c_k \frac{\lambda^k}{k!}\right).$$

In the noncommutative setting, these are replaced by ordinary generating functions.

Let

$$M(w) \coloneqq 1 + \sum_{\alpha} \varphi(a_{\alpha}) w_{\alpha}, \quad R(w) \coloneqq \sum_{\alpha} \kappa(a_{\alpha}) w_{\alpha}, \quad \eta(w) \coloneqq \sum_{\alpha} \beta(a_{\alpha}) w_{\alpha}.$$

Considering a new set of variables $z_i = w_i M(w)$ we have

$$M(w)=1+R(z),\quad M(w)=1+\eta(w)M(w).$$



It turns out that the Hopf-algebraic language above describes two operations on power series.

Let G^{p} and G^{c} denote the group of invertible power series and formal diffeomorphisms, resp.

For $f, g \in G^p$ define

$$f^g(w) \coloneqq g(w)f(z), \quad z_i = w_i g(w).$$

Also let

$$(f \curvearrowleft g)(w) \coloneqq f(z), \quad z_i = w_i g(w).$$

Given $F : T(A) \rightarrow k$ let $\Lambda(F) \in k[[w]]$ be given by

$$\Lambda(F)(w) = F(1) + \sum_{\alpha} F(a_{\alpha})w_{\alpha}.$$



Theorem

$$\Lambda(F * G) = \Lambda(F)^{\Lambda(G)}$$

Theorem

 $\Lambda(F \prec G) = \Lambda(F) \curvearrowleft \Lambda(G).$

Theorem

 $\Lambda(F \succ G) = \Lambda(F) \frown \Lambda(G).$

